

# What is DNA Sequencing?

Genomes are composed of **base pairs**: Adenine, Guanine, Cytosine, Thymine.

The goal is to find the **sequence** of base pairs which compose the genome.

### Why is this useful?

- Tracing evolution.
- Correlating genes with diseases.
- Forensics and identification.



# **How is Sequencing Done?**

- DNA split into small pieces called **reads**.
- Using a **reference** genome, reads are mapped to potential locations.
- Must account for errors in reads: insertions, deletions, and **substitutions**.



This problem is very **computationally challenging**:

- Billions of reads.
- Fuzzy string matching.
- Multiple mapping locations per read.

# **Research Question**

**Question:** Can we exploit multi-core machines to improve speed?

- Offset high cost of computation: read mapping is slow.
- Efficiently use resources (caches, cores, etc).
- Achieve good speedup and minimize overhead of concurrency.
- Build a low complexity, cache-efficient, memory efficient system.
- DNA Infrastructure is inherently parallel.

### **Machine Specifications:**

- 4 sockets.
- 10 cores and 256GB RAM (NUMA) per socket.
- Intel Hyperthreading.
- Total 80 logical threads, 1TB RAM.

# Parallelizing Seed Selection in DNA Read Mapping Sunny Nahar

GCTATGGTAATCGCTGCATATC genome length  $G \approx 10^9$  $N \approx 10^8$ GGACACCCAAGG GGACACCCAAGGGGGAA

# Generating the HashTree

### **Description**:

- Hashtable of frequency **tries** of the reference genome which stores seed frequencies.
- Queries are **O(L)** cache misses.
- Need to load from disk to construct and perform computation: 80GB.

### **Parallelization**:

- Each thread reads portion of file.
- Using memory mapped IO removes copy to main memory (kernel page cache).
- Each trie can be independently generated. Construction time hides disk latencies.
- Dynamic work scheduling.



# **Bidirectional Frequency Predictor Construction**

Tool for speeding up seed selection. A data structure to **estimate** frequency of seeds:

- Used to **minimize** reference trie queries.
- Predicts frequency given base seed, left and right extension.
- Gives **O(1)** cache miss complexity.

### **Parallelization:**

- Requires single DFS traversal of HashTree.
- Threads take set of tries from the hashtable.
- Only synchronization is atomic writes.





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5	0	60	70	8

# **Seed Selection:**



## **Parallel Seed Selection** Given a read, output set of seeds. Used in next stage (edit distance filtering). • Low frequency seeds are important. ATCTACGAGGCTACTAGCTAGGCATCTACTATCTACTACAGCGACGGCGG 492 620 439 503 Freq: There are multiple seed selection algorithms and heuristics:

- Mix of accesses to HashTree and predictor.
- Varying levels of complexity.

## **Parallelization**:

- Parallelize selection over read set.
- HashTree and predictor are static: read coherence overhead.
- NUMA effects on non-local threads (up to 2x overhead of read).



# **Conclusion and Future Work**

the seed selection pipeline.

### Exhibit large speedups:

- 15x for HashTree construction.
- 41x for Frequency Predictor construction.
- 60-70x for Seed Selection algorithms.
- Still room for further optimizations.
  - Processor affinity. • Masking NUMA with work distribution.

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We were able to efficiently parallelize various components of