Parallelizing DNA Read Mapping

 $\bullet \bullet \bullet$

Sunny Nahar

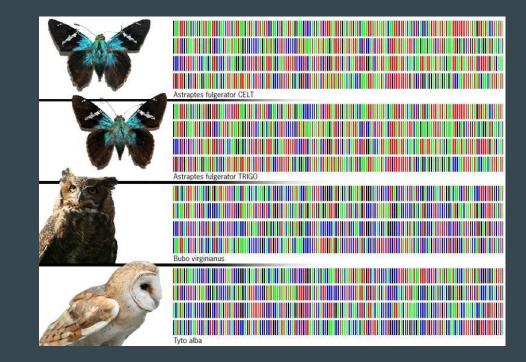
What is DNA Sequencing?

Finding the base-pairs for the genome.



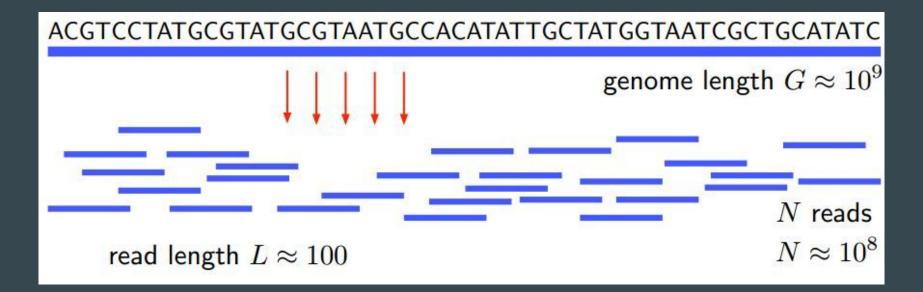
Why is this useful?

- Tracing evolution.
- Correlating genes with diseases.
- Forensics and identification.



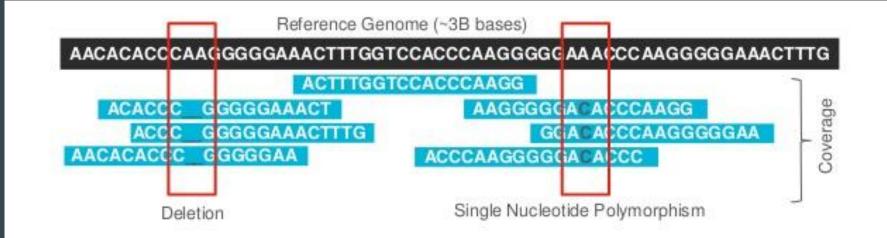
Current Technology (Shotgun)

• Split DNA into small pieces (reads).



Read Mapping

- Have access to a reference genome.
- Align reads to reference.

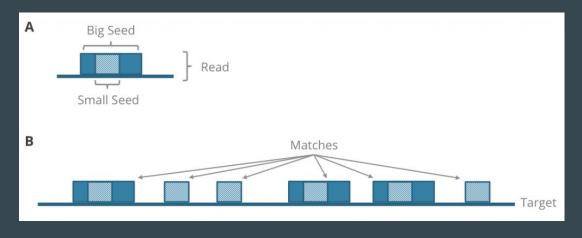


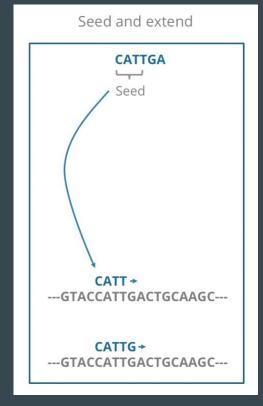
Computationally Challenging

- Billions of reads.
- Fuzzy matching.
 - Handle insertions, deletions, mutations, errors.
- Multiple mapping locations.
- Assemble with high probability.

How do we map a read? (Seed-and-extend Method)

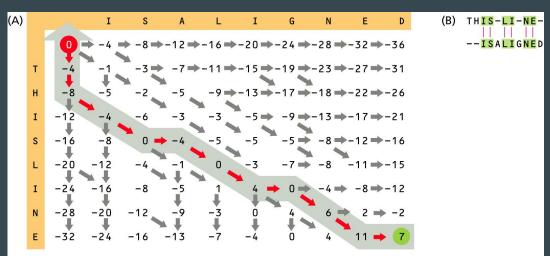
- Match substrings (seeds) exactly to the reference.
 - Possible locations.





How do we map a read? (Seed-and-extend Method)

- Use edit distance to determine quality.
 - Dynamic Programming (score based)
 - Needleman-Wunsch
 - Smith-Waterman
- Choosing less frequent substrings is important.



Research Project

- Improve the speed of the mapper (aligning reads).
- Develop novel algorithms and heuristics.
- Low complexity, memory efficient, cache efficient.

This Presentation

Focuses on one part of the pipeline: Seed Selection.
Given set of reads, output seeds.
These are used to find potential mapping locations.
Discuss parallel optimizations and improvements to runtime.

Parallelizing the Infrastructure

DNA Read Processing Pipeline

- 1. Generating the HashTree representation of genome.
- 2. Building a frequency predictor.
- 3. Performing seed selection.
- 4. Pipe results into next stage (edit distance).

Machine Specs

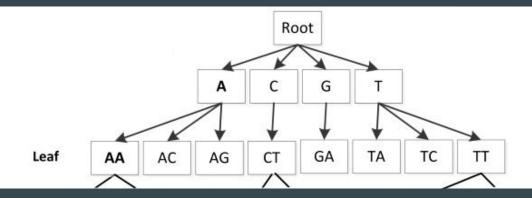
Test machine:

- 4 sockets.
- 10 cores and 256GB RAM (NUMA) per socket.
- 2 hardware threads per core (Intel Hyperthreading).
 Memory latency.
- In total, 1TB RAM with 80 logical threads.

Generating the HashTree

Genome Representation

- Hashtable of Frequency Tries.
 - Each node stores a character and frequency.
 - \circ $\;$ Hashtable on first 10 letters.
- Bounded (length 30)
- ~80GB on disk.
- String frequency queries.
- L cache misses.



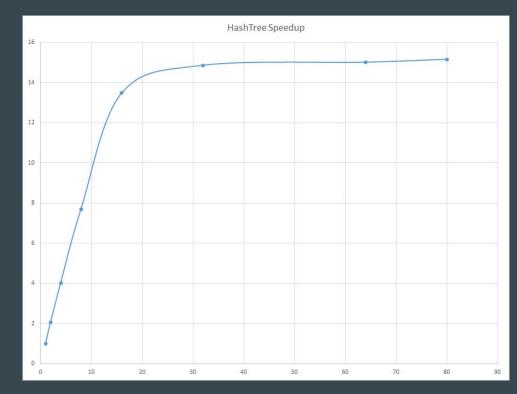
Generation in Parallel

- Reading from disk is sequential.
 - Threads take turn reading.
 - Memory mapped IO removes need for explicit copying.
 - Copied to Kernel page cache as opposed to user memory.
 - Incur TLB misses vs cache misses.
- Each trie can be independently generated.
 - Only issue is memory allocator.
 - Traditional malloc has a lock.
 - Only need alloc (not free).
 - Implement own allocator which is locality aware.
 - Was initial bottleneck. (2hrs to 10 minutes)

Generation in Parallel

- Dynamic work scheduling + greedy allocation.
 - Trie sizes are highly nonuniform.
 - Schedule largest tries first to balance workload.
 - Estimate from filesize.
- Kernel aware access policies.
 - TLB linear access.
 - File linear access.

Speedup Graph



Frequency Predictor

 $\bullet \bullet \bullet$

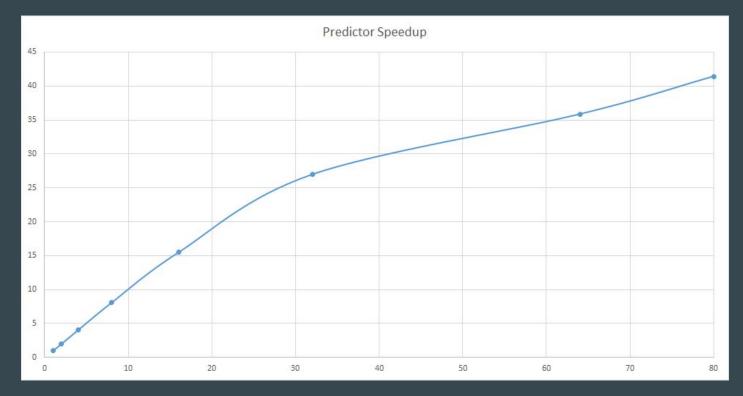
What is the Frequency Predictor?

- Access to HashTree is costly (L cache misses)
 - Instead, give an estimated frequency.
- Reduces to 1 cache miss.
- Store a table:
 - table[base][L][R] -> base (10 letters) extended to left by L letters, right by R letters.
- Example: AGCTGACG ATGCTAGCTA GCTCG
 Lookup table[ATGCTAGCTA][8][5]

Construction of Predictor

- Requires traversing through the entire HashTree.
- Updating a large table
 - Synchronization at same entries.
 - Accomplished with atomic writes.

Speedup Graph



Seed Selection

 $\bullet \bullet \bullet$

What is Seed Selection?

- Given input set of reads, output set of seeds for each read.
 - Based on input parameters.
 - GCAGTCAGTCGATCGATCGATCGTACGTACGTACAGCTAGC TA
- Algorithms use mix of accesses to HashTree and predictor to determine seeds.

Parallelization

- Selection is parallelized over reads.
 - Per thread data structures, reduced at end.
- Both the HashTree and Predictor are loaded in memory.
 - Generally sparse accesses.
 - Cache write coherence is not an issue, since only read access to memory.
 - Cache reads create coherence traffic (costly on socket architecture).

Parallelization

- Stack memory vs malloc.
- NUMA degrades performance.
 - Threads closest to HashTree, Predictor perform much faster.
 - Observed up to 2x overhead.

Speedup Graph



Questions

 $\bullet \bullet \bullet$